# Descriptive set theory and geometrical paradoxes II 

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## Borel circle squaring

## Theorem (M.-Unger, 2016)

Tarski's circle squaring problem can be solved using Borel pieces. More generally, suppose $k \geq 1$ and $A, B \subseteq \mathbb{R}^{k}$ are bounded Borel sets such that $\lambda(A)=\lambda(B)>0, \Delta(\partial A)<k$, and $\Delta(\partial B)<k$. Then $A$ and $B$ are equidecomposable by translations using Borel pieces.
$\lambda$ is Lebesgue measure, and $\Delta$ is upper Minkowski dimension.

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$\lambda$ is Lebesgue measure, and $\Delta$ is upper Minkowski dimension.
Fix $k$ and such sets $A$ and $B$.

## Laczkovich's ideas: Work in the torus

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Fix a sufficiently large $d$ and randomly pick $u_{1}, \ldots, u_{d} \in \mathbb{T}^{k}$. Obtain an action $a$ of $\mathbb{Z}^{d}$ on $\mathbb{T}^{k}$ by letting the $i$ th generator of $\mathbb{Z}^{d}$ act via $u_{i}$.

$$
\left(n_{1}, \ldots, n_{d}\right) \cdot x=n_{1} u_{1}+\ldots+n_{d} u_{d}+x
$$

Laczkovich shows $A$ and $B$ are a-equidecomposable.

## Laczkovich's ideas: Discrepancy theory

If $F \subseteq \mathbb{T}^{k}$ is finite and $C \subseteq \mathbb{T}^{k}$ is Lebesgue measurable, then the discrepancy of $F$ with respect to $C$ is

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## Lemma (Laczkovich 1992 building on Schmidt, Niederreiter-Wills)

For $A, B$ and the action as above, $\exists \epsilon>0$ and $M$ such that

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D\left(R_{N} \cdot x, A\right) \leq M N^{-1-\epsilon} \text { and } D\left(R_{N} \cdot x, B\right) \leq M N^{-1-\epsilon} .
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Roughly, every square of side length $N$ in the action contains close to $\lambda(A) N^{d}$ elements of both $A$ and $B$.

## Flows in graphs

Suppose $G$ is a graph (symmetric irreflexive relation) on a vertex set $X$. If $f: X \rightarrow \mathbb{R}$ is a function, then an $f$-flow of $G$ is a function $\phi: G \rightarrow \mathbb{R}$ such that

- For every edge $(x, y) \in G, \phi(x, y)=-\phi(y, x)$, and
- For every vertex $x \in X$,

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## Flows and equidecompositions

For the rest of the proof, let $G$ be the graph with vertex set $\mathbb{T}^{k}$ where $x, y \in \mathbb{T}^{k}$ are adjacent if there is $g \in \mathbb{Z}^{d}$ such that $g \cdot x=y$ where $|g|_{\infty}=1$.

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## Proposition

$A$ and $B$ are a-equidecomposable with Borel pieces iff there is a bounded Borel integer-valued $\chi_{A}-\chi_{B}$-flow of $G$.
$\rightarrow: A$ and $B$ are a-equidecomposable with Borel pieces iff there is Borel bijection $\theta: A \rightarrow B$ and a finite set $S \subseteq \mathbb{Z}^{d}$ such that $\forall x \in A \exists g \in S(\theta(x)=g \cdot x)$.

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To construct a flow, for each $x \in A$ add 1 unit of flow to each edge along the lex-least path from $x$ to $\theta(x)$.

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Find a Borel tiling of each orbit by rectangles of side length $\approx N$. Each tile has roughly $\lambda(A) N^{d}$ points of $A$ and $B$, and the flow over the boundary of the tile is $\leq O\left(c N^{d-1}\right)$. Using discrepancy, if $N$ is sufficiently large, there are more points of $A$ and $B$ in every tile than maximum flow out of the boundary of the tile.

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Now construct a Borel bijection from $A$ to $B$ witnessing equidecomposability. Suppose $R, S$ are adjacent tiles. If

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\sum_{(x, y) \in G: x \in R \wedge y \in S} \phi(x, y)>0
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map this many points of $A \in R$ to points of $B \in S$. If the quantify is negative, map this many points of $B \in R$ to $A \in S$.

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map this many points of $A \in R$ to points of $B \in S$. If the quantify is negative, map this many points of $B \in R$ to $A \in S$. Since $\phi$ is a $\chi_{A}-\chi_{B}$-flow, after doing this the same number of points of $A$ and $B$ remain in each tile. Biject them to finish the construction.

## How to construct tilings

An independent set in a graph $G$ is a set of vertices where no two are adjacent.

Theorem (Kechris, Solecki, Todorcevic, 1999)
If $G$ is a locally finite Borel graph, then there is a Borel maximal independent set for $G$.

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Let $G \leq n$ be the graph on $\mathbb{T}^{k}$ where $x, y$ are adjacent if $d_{G}(x, y) \leq n$. Let $C$ be a Borel maximal independent set for $G \leq n$. Use the element of $C$ as center points for "tiles" of $G$.

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Let $G^{\leq n}$ be the graph on $\mathbb{T}^{k}$ where $x, y$ are adjacent if $d_{G}(x, y) \leq n$. Let $C$ be a Borel maximal independent set for $G \leq n$. Use the element of $C$ as center points for "tiles" of $G$.

If we use these center points to make "Voroni cells", the resulting tiling suffices. Gao-Jackson (2015) give a more complicated construction to make rectangular tilings.

## A sketch of our proof

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2. We show that given any real-valued Borel $f$-flow of $G$, we can find an integer valued Borel $f$-flow which is "close" to the real-valued one.
3. We finish by using the proposition we've proved above: there's a Borel equidecomposition iff there is a bounded Borel $\chi_{A}-\chi_{B}$-flow.

Finding a real-valued bounded flow

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For every $i>0$, let $\pi_{i}: \mathbb{Z}^{d} /\left(2^{i} \mathbb{Z}\right)^{d} \rightarrow \mathbb{Z}^{d} /\left(2^{i-1} \mathbb{Z}\right)^{d}$ be the canonical homomorphism. This yields the inverse limit
where elements of $\hat{\mathbb{Z}^{d}}$ are sequences $\left(h_{0}, h_{1}, \ldots\right)$ such that $\pi_{i}\left(h_{i}\right)=h_{i-1}$ for all $i>0$.

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For each $h \in \hat{\mathbb{Z}}^{d}$ and $x \in \mathbb{T}^{k}$, we give an explicit construction $\phi_{x, h}$ of a flow of the connected component of $x$. However, we cannot pick a single $x$ in each orbit to be a "starting point" for this construction (since this would be a nonmeasurable Vitali set).

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The construction is such that if $g \in \mathbb{Z}^{d}$, then $\phi_{x, h}=\phi_{g \cdot x,-g+h}$. Hence, the average value of this construction is invariant of our starting point ( $h \mapsto-g+h$ is measure preserving):

$$
\int_{h} \phi_{x, h}=\int_{h} \phi_{g \cdot x,-g+h}=\int_{h} \phi_{g \cdot x, h}
$$

In the last lecture, we'll discuss how to turn a real-valued flow of $G$ into an integer-valued flow. This step uses:

- the Ford-Fulkerson algorithm in finite combinatorics.
- work of $A$. Timár on boundaries of finite sets in $\mathbb{Z}^{d}$.
- very recent work of Gao, Jackson, Krohne and Seward on hyperfiniteness of free Borel actions of $\mathbb{Z}^{d}$.

